## Ciências ULisboa

## Percolation

Nuno Araújo
Centro de Física Teórica e Computacional, Universidade de Lisboa, Portugal

## Books on percolation

$\rightarrow$ D. Stauffer and A. Aharony, Introduction to percolation theory. CRC Press (2000).
$>$ M. Sahimi, Applications of percolation theory. Taylor \&
Francis (1994).
$>$ K. Christensen and N. R. Moloney, Complexity and criticality. Imperial College Press (2005).

## Forest fire



## Forest fire



## Spreading of epidemics



## Spreading of epidemics



## Oil fields


at Barrancabermeja (Colombia), photo by Melissa Jiménez.

## Percolation model



## Percolation model



Percolation model
order parameter

$$
P_{\infty}=\frac{S_{\max }}{N}
$$



$$
P_{\infty} \sim\left(p-p_{c}\right)^{\beta}
$$

## Percolation threshold

## largest cluster: fractal dimension



## Percolation threshold cluster-size distribution

$$
n_{s} \sim s^{-\tau}
$$




## Algorithms

generate canonical configurations


## For each site $i$ :

1. random number $\varepsilon$;
2. if
$\boldsymbol{\varepsilon}<\mathbf{p}$ : $\boldsymbol{i}$ is occupied;
else: $\boldsymbol{i}$ is empty.

## Algorithms <br> Burning method



1. set first row burning;
H. J. Herrmann, D. C. Hong, and H. E. Stanley. J. Phys. A 17, L261 (1984)

## Algorithms

1. set first row burning;
2. set neighbors of burning to burning and burning to burned;

## Algorithms <br> Burning method

burning
occupied
burned


1. set first row burning;
2. set neighbors of burning to burning and burning to burned;
3. repeat until everything is burned.

## Algorithms <br> Burning method

burning
occupied
burned


1. set first row burning;
2. set neighbors of burning to burning and burning to burned;
3. repeat until everything is burned.

## Algorithms



1. set first row burning;
2. set neighbors of burning to burning and burning to burned;
3. repeat until everything is burned.

## Algorithms



1. set first row burning;
2. set neighbors of burning to burning and burning to burned;
3. repeat until everything is burned.

## Algorithms



1. set first row burning;
2. set neighbors of burning to burning and burning to burned;
3. repeat until everything is burned.

## Algorithms



1. set first row burning;
2. set neighbors of burning to burning and burning to burned;
3. repeat until everything is burned.

## Algorithms

1. set first row burning;
2. set neighbors of burning to burning and burning to burned;
3. repeat until everything is burned.

## Algorithms <br> Burning method

$\bigcirc$
empty
burning
occupied
burned


1. set first row burning;
2. set neighbors of burning to burning and burning to burned;
3. repeat until everything is burned.

## Algorithms <br> Burning method

empty
burning
occupied
burned


## One can determine if the set of occupied sites percolates or not.



Number of clusters and cluster size distribution?

## Algorithms

$$
k=2
$$

## Hoshen and Kopelman

$\mathrm{M}(\mathrm{k})=0$


1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.
J. Hoshen and R. Kopelman. Phys. Rev. B 14, 3438 (1976)

## Algorithms

## Hoshen and Kopelman



1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 0 |

## Algorithms

## Hoshen and Kopelman



1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 1 |

## Algorithms

## Hoshen and Kopelman



1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 2 |

## Algorithms

## Hoshen and Kopelman

1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and
 bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 2 |
| 3 | 1 |

## Algorithms

## Hoshen and Kopelman

1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and
 bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 1 |

## Algorithms

## Hoshen and Kopelman

1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and
 bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 1 |
| 4 | 1 |

## Algorithms

## Hoshen and Kopelman

1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and
 bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 3 |
| 4 | -3 |

## Algorithms

## Hoshen and Kopelman

1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and
 bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 3 |
| 4 | -3 |
| 5 | 1 |
|  |  |

## Algorithms

## Hoshen and Kopelman

1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and
 bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 4 |
| 4 | -3 |
| 5 | 1 |
|  |  |

## Algorithms

## Hoshen and Kopelman

1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and
 bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 7 |
| 4 | -3 |
| 5 | 1 |

## Algorithms

## Hoshen and Kopelman

1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and
 bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 8 |
| 4 | -3 |
| 5 | 1 |
| 6 | 1 |

## Algorithms

## Hoshen and Kopelman

1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and
 bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 8 |
| 4 | -3 |
| 5 | 1 |
| 6 | 2 |

## Algorithms

## Hoshen and Kopelman



1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 8 |
| 4 | -3 |
| 5 | 1 |
| 6 | 2 |
| 7 | 1 |

## Algorithms

## Hoshen and Kopelman



1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 8 |
| 4 | -3 |
| 5 | 1 |
| 6 | 2 |
| 7 | 2 |

## Algorithms

## Hoshen and Kopelman



1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 8 |
| 4 | -3 |
| 5 | 1 |
| 6 | 2 |
| 7 | 2 |
| 8 | 1 |

## Algorithms

## Hoshen and Kopelman



1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 10 |
| 4 | -3 |
| 5 | 1 |
| 6 | 2 |
| 7 | 2 |
| 8 | -3 |

## Algorithms

## Hoshen and Kopelman



1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 11 |
| 4 | -3 |
| 5 | 1 |
| 6 | 2 |
| 7 | 2 |
| 8 | -3 |

## Algorithms

## Hoshen and Kopelman



1. start from the site in the left-bottom corner;
2. sweep from left to right bottom to top;
3. only verify left and bottom neighbors.

| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 11 |
| 4 | -3 |
| 5 | 1 |
| 6 | 2 |
| 7 | 2 |
| 8 | -3 |

## Algorithms

## Hoshen and Kopelman



| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 11 |
| 4 | -3 |
| 5 | 1 |
| 6 | 2 |
| 7 | 2 |
| 8 | -3 |
| 9 | 1 |

J. Hoshen and R. Kopelman. Phys. Rev. B 14, 3438 (1976)

## Algorithms

## Hoshen and Kopelman




Two neighbor ko:
$M\left(\underline{k_{0}}\right)=M\left(\underline{k_{0}}\right)+1$

One neighbor ko:
$M\left(\underline{k_{0}}\right)=M\left(\underline{k_{0}}\right)+1$

J. Hoshen and R. Kopelman. Phys. Rev. B 14, 3438 (1976)

## Algorithms

Newman and Ziff (microcanonical)


| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 0 |

M. E. J. Newman and R. M. Ziff. Phys. Rev. Lett. 85, 4104 (2000) M. E. J. Newman and R. M. Ziff. Phys. Rev. E 64, 016706 (2001)

## Algorithms

Newman and Ziff (microcanonical)


| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 1 |

M. E. J. Newman and R. M. Ziff. Phys. Rev. Lett. 85, 4104 (2000) M. E. J. Newman and R. M. Ziff. Phys. Rev. E 64, 016706 (2001)

## Algorithms

## Newman and Ziff (microcanonical)



| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 1 |
| 3 | 1 |

M. E. J. Newman and R. M. Ziff. Phys. Rev. Lett. 85, 4104 (2000) M. E. J. Newman and R. M. Ziff. Phys. Rev. E 64, 016706 (2001)

## Algorithms

## Newman and Ziff (microcanonical)



| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 1 |
| 3 | 2 |

M. E. J. Newman and R. M. Ziff. Phys. Rev. Lett. 85, 4104 (2000) M. E. J. Newman and R. M. Ziff. Phys. Rev. E 64, 016706 (2001)

## Algorithms

## Newman and Ziff (microcanonical)



| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 1 |
| 3 | 2 |
| 4 | 1 |

M. E. J. Newman and R. M. Ziff. Phys. Rev. Lett. 85, 4104 (2000) M. E. J. Newman and R. M. Ziff. Phys. Rev. E 64, 016706 (2001)

## Algorithms

## Newman and Ziff (microcanonical)



| $k$ | $M(k)$ |
| :---: | :---: |
| 2 | 3 |
| 3 | 2 |
| 4 | -2 |

M. E. J. Newman and R. M. Ziff. Phys. Rev. Lett. 85, 4104 (2000) M. E. J. Newman and R. M. Ziff. Phys. Rev. E 64, 016706 (2001)

## Algorithms

## Microcanonical vs canonical



Fixed number of
occupied sites (n)

Fixed probability that a site is occupied ( $p$ )

$$
\boldsymbol{B}(\boldsymbol{N}, \boldsymbol{n}, \boldsymbol{p})=\binom{N}{n} p^{n}(1-p)^{N-n}
$$

$B(N, n, p)$ : probability that exactly n sites are occupied in a canonical configuration

$$
Q(p)=\sum_{n=0}^{N} \boldsymbol{B}(\boldsymbol{N}, \boldsymbol{n}, \boldsymbol{p}) Q_{n}=\sum_{n=0}^{N}\binom{N}{n} p^{n}(1-p)^{N-n} Q_{n}
$$

M. E. J. Newman and R. M. Ziff. Phys. Rev. Lett. 85, 4104 (2000) M. E. J. Newman and R. M. Ziff. Phys. Rev. E 64, 016706 (2001)

